



Polynomial Multiplication Techniques

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June 8, 2023 at Vodice

Polynomial Multiplications

Karatsuba and Toom(-Cook)



Our Goal

- Do multiplication over rings like $\mathbb{Z}_{3329}[x]/\langle x^{256} + 1 \rangle$ or $\mathbb{Z}_{4591}[x]/\langle x^{761} - x - 1 \rangle$



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- For efficiency considerations, one would like to replace multiplication sub-steps by addition sub-steps as much as possible
- This talk reviews some techniques for this purpose



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- Do redundant modular arithmetic

e.g. $f(x), g(x) \in \mathbb{Z}_p[x]$ with $\deg(f) + \deg(g) < n \longrightarrow \bar{f}(x), \bar{g}(x) \in \mathbb{Z}_p[x]/\langle x^n - 1 \rangle$



Karatsuba

- A simple observation: $(a_0 + a_1x)(b_0 + b_1x) = a_0b_0 + (a_0b_1 + a_1b_0)x + a_1b_1x^2$
where $a_0b_1 + a_1b_0 = (a_0 + a_1)(b_0 + b_1) - a_0b_0 - a_1b_1$

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- Improved form: $(a_0 + a_1x)(b_0 + b_1x) = (a_0 + a_1)(b_0 + b_1)x + (a_0b_0 - a_1b_1x)(1 - x)$



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- This applies when our operands are in the form of linear polynomials



Karatsuba to multiply $f(x), g(x)$ each having degree $< 2n$

- Change x^n to y , use Karatsuba in y : 3 polynomials in x of degree $< n$

$$(1 + 2x + 2x^2 + 2x^3) \cdot (3 + x + 4x^2 + x^3) = [(1 + 2x) + (2 + 2x)y] \cdot [(3 + x) + (4 + x)y]$$



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$$\begin{aligned} & (1 + 2x + 2x^2 + 2x^3) \cdot (3 + x + 4x^2 + x^3) = [(1 + 2x) + (2 + 2x)y] \cdot [(3 + x) + (4 + x)y] \\ = & \quad [(3 + 4x)(7 + 2x)]y \quad + [(1 + 2x)(3 + x) - (2 + 2x)(4 + x)y](1 - y) \\ = & [63x + (21 - 8x)(1 - x)]y \quad + [12x + (3 - 2x)(1 - x) - [20x + (8 - 2x)(1 - x)]y](1 - y) \\ = & [21 + \binom{-8+63}{-21}x + 8x^2]y \quad + [3 + \binom{-2+12}{-3}x + 2x^2 - [8 + \binom{-2+20}{-8}x + 2x^2]y](1 - y) \\ = & [21 + 34x + 8x^2]y \quad + [3 + 7x + 2x^2 - (8 + 10x + 2x^2)y](1 - y) \end{aligned}$$



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Why “Improved” Karatsuba? i

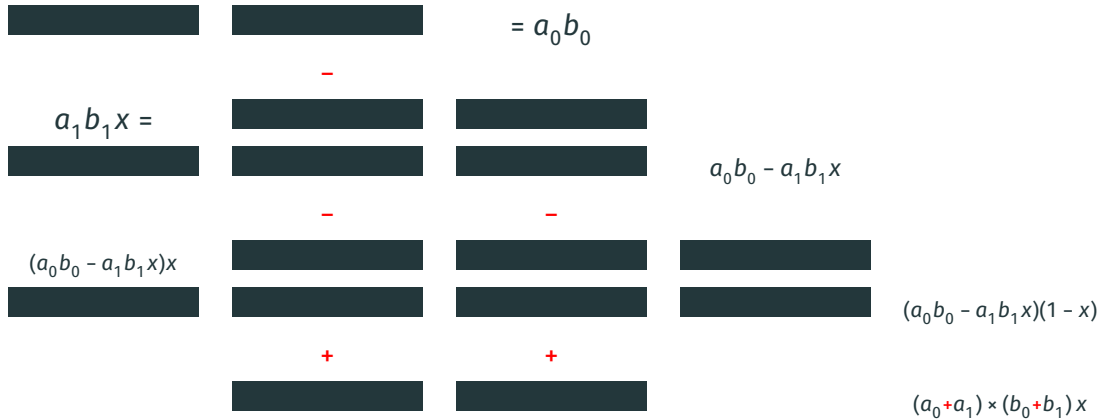
Counting Additions: Why is $(a_0 + a_1x)(b_0 + b_1x) = (a_0 + a_1)(b_0 + b_1)x + (a_0b_0 - a_1b_1x)(1 - x)$ better?

- Suppose $x = t^{100}$, each of a_0, a_1, b_0, b_1 is a length 100 polynomial in t .
- Each product $a_0b_0, a_1b_1, (a_0 + a_1)(b_0 + b_1)$ is a length 199 polynomial in t .
- Can count 8 additions/subtractions in “standard” Karatsuba



Why “Improved” Karatsuba? ii

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(actually 100) addition/subtraction has seemingly vanished into thin air!!

Repeated Karatsuba: Memory Access

If we want to do three layers of Karatsuba for polynomials of degree $< 8n$

■: polynomial of degree $< n$

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Toom-3

- To multiply $(a_0 + a_1x + a_2x^2)(b_0 + b_1x + b_2x^2)$, evaluate at 5 points
A simple choice will be $x = 0, \pm 1, -2, \infty \rightarrow F(0), F(1), F(-1), F(-2), F(\infty)$



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A simple choice will be $x = 0, \pm 1, -2, \infty \rightarrow F(0), F(1), F(-1), F(-2), F(\infty)$
- Interpolate the degree-4 polynomial $F(x) = \sum_{i=0}^4 c_i x^i$. In matrix form,

$$\begin{bmatrix} F(0) \\ F(1) \\ F(-1) \\ F(-2) \\ F(\infty) \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 & 1 \\ 1 & -2 & 4 & -8 & 16 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \\ c_4 \end{bmatrix}$$

The coefficients can be determined by applying the inverse matrix

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The coefficients can be determined by applying the inverse matrix

- This applies when our operands are in the form of polynomials of degree < 3

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- To multiply $f(x), g(x)$ each having degree $< 4n$, change x^n to y



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- Use Toom-4 in y : evaluate y at 7 points (= 7 mult. of poly. in x of degree $< n$)



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- To multiply $f(x), g(x)$ each having degree $< 4n$, change x^n to y
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- Use Toom-4 in y : evaluate y at 7 points (= 7 mult. of poly. in x of degree $< n$)
- Interpolate the polynomial in y
- Sometimes uses plus-minus powers of 2, not integers as interpolation points.

“Evaluation and Interpolation at $1/a$ ”

Instead of computing $f(1/a) = \sum_{j=0}^{k-1} f_j(1/a)^j$, we compute

$a^{k-1} f(1/a) = \sum_{j=0}^{k-1} f_j a^{k-1-j}$. After the point multiplication, we have

$$a^{2k-2} f(1/a)g(1/a) = \left(\sum_{j=0}^{k-1} f_j a^{k-1-j} \right) \left(\sum_{j=0}^{k-1} g_j a^{k-1-j} \right).$$



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 - May be best for NTRU type on Neon with Toeplitz variation.
 - Note you need extra precision to divide by $2 \bmod 2^k$.



Toeplitz Matrix Methods

Want these matrix-vector product, which cyclic convolutions are

$$\begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \\ \vdots \\ c_{n-1} \end{bmatrix} = \begin{bmatrix} a_0 & a_{-1} & a_{-2} & a_{-3} & \cdots & a_{-n+1} \\ a_1 & a_0 & a_{-1} & a_{-2} & \cdots & a_{-n+2} \\ a_2 & a_1 & a_0 & a_{-1} & \cdots & a_{-n+3} \\ a_3 & a_2 & a_1 & a_0 & \cdots & a_{-n+4} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{n-1} & a_{n-2} & a_{n-3} & a_{n-4} & \cdots & a_0 \end{bmatrix} \begin{bmatrix} b_0 \\ b_1 \\ b_2 \\ b_3 \\ \vdots \\ b_{n-1} \end{bmatrix}$$

Such matrices are “Toeplitz”. Submatrices of a Toeplitz matrix are Toeplitz. So

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} A_0 & A_{-1} \\ A_1 & A_0 \end{bmatrix} \begin{bmatrix} B_0 \\ B_1 \end{bmatrix} \text{ or } \begin{cases} c_0 = A_{-1}B_1 + A_0B_0 = (A_{-1} - A_0)B_1 + A_0(B_0 + B_1), \\ c_1 = A_0B_1 + A_1B_0 = A_0(B_0 + B_1) + (A_1 - A_0)B_0. \end{cases}$$

How to Obtain Toeplitz formulas from Toom/Karatsuba formulas

$$B_0 C_0 = B_0 C_0$$

$$B_0 C_1 + B_1 C_0 = (B_0 + B_1)(C_0 + C_1) - B_0 C_0 - B_1 C_1$$

$$B_1 C_1 = B_1 C_1$$

Now multiply the three formulas by A_0 , A_1 , and A_2 , add together

$$A_0 B_0 C_0 + A_1 B_0 C_1 + A_1 B_1 C_0 + A_2 B_1 C_1 = (A_0 - A_1) B_0 C_0 + A_1 (B_0 + B_1) (C_0 + C_1) + (A_2 - A_1) B_1 C_1$$

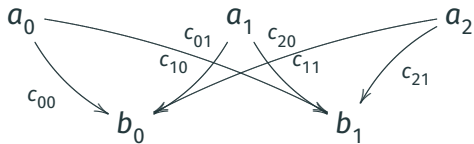
Now collect the terms according to the C_i :

$$C_0 : A_0 B_0 + A_1 B_1 = (A_0 - A_1) B_0 + A_1 (B_0 + B_1)$$

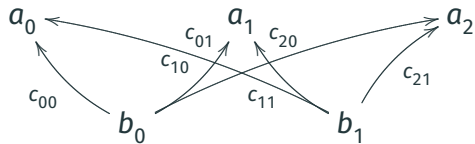
$$C_1 : A_1 B_0 + A_2 B_1 = A_1 (B_0 + B_1) + (A_2 - A_1) B_1$$

Transposition of Linear Maps

A tagged arrow means to multiply by tag and add to target



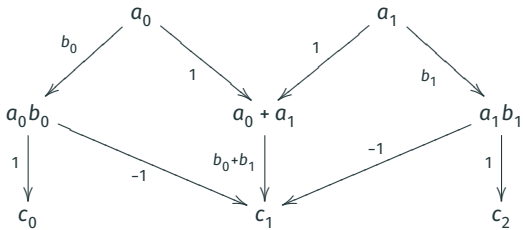
$$\begin{bmatrix} b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} c_{00} & c_{10} & c_{20} \\ c_{01} & c_{11} & c_{21} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \end{bmatrix}$$



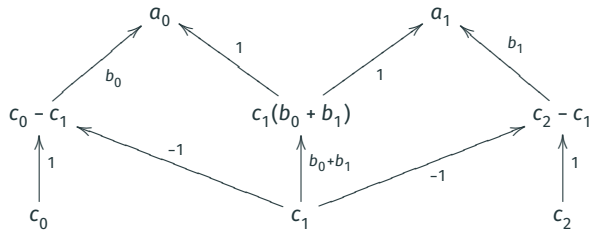
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Transposition of Karatsuba

Transposition of usual polynomial product is Toeplitz-Matrix-to-Vector Product



$$\begin{aligned} c_0 &= a_0 b_0 \\ c_1 &= (a_0 + a_1)(b_0 + b_1) - a_0 b_0 - a_1 b_1 \\ c_2 &= a_1 b_1 \end{aligned}$$



$$\begin{aligned} a_0 &= (c_0 - c_1)b_0 + c_1(b_0 + b_1) \\ a_1 &= c_1(b_0 + b_1) + (c_2 - c_1)b_1 \end{aligned}$$

TMVP formulations for NTRU variants

$c = ab$ in NTRU Ring $\mathbb{Z}_q[x]/(x^p - 1)$ as TMVP

$$\begin{aligned}
 c_0 &= a_0 b_0 + a_1 b_{p-1} + a_2 b_{p-2} + \dots + a_{p-2} b_2 + a_{p-1} b_1 \\
 c_1 &= a_0 b_1 + a_1 b_0 + a_2 b_1 + \dots + a_{p-2} b_3 + a_{p-1} b_2 \\
 c_2 &= a_0 b_2 + a_1 b_1 + a_2 b_0 + \dots + a_{p-2} b_4 + a_{p-1} b_3 \\
 &\vdots \\
 c_{p-2} &= a_0 b_{p-2} + a_1 b_{p-3} + a_2 b_{p-4} + \dots + a_{p-2} b_0 + a_{p-1} b_{p-1} \\
 c_{p-1} &= a_0 b_{p-1} + a_1 b_{p-2} + a_2 b_{p-3} + \dots + a_1 b_{p-2} + a_{p-1} b_0
 \end{aligned}$$

$c = ab$ in NTRU Prime Ring $\mathbb{Z}_q[x]/(x^p - x - 1)$ as TMVP

$$\begin{aligned}
 c_0 + c_{p-1} &= a_0(b_0 + b_{p-1}) + a_1(b_{p-1} + b_{p-2}) + a_2(b_{p-2} + b_{p-3}) + \dots + a_{p-2}(b_2 + b_1) + a_{p-1}(b_1 + b_0 + b_{p-1}) \\
 c_1 &= a_0 b_1 + a_1(b_0 + b_{p-1}) + a_2(b_1 + b_{p-2}) + \dots + a_{p-2}(b_3 + b_2) + a_{p-1}(b_2 + b_1) \\
 c_2 &= a_0 b_2 + a_1 b_1 + a_2(b_0 + b_{p-1}) + \dots + a_{p-2}(b_4 + b_3) + a_{p-1}(b_3 + b_2) \\
 &\vdots \\
 c_{p-2} &= a_0 b_{p-2} + a_1 b_{p-3} + a_2 b_{p-4} + \dots + a_{p-2}(b_0 + b_{p-1}) + a_{p-1}(b_{p-1} + b_{p-2}) \\
 c_{p-1} &= a_0 b_{p-1} + a_1 b_{p-2} + a_2 b_{p-3} + \dots + a_1 b_{p-2} + a_{p-1}(b_0 + b_{p-1})
 \end{aligned}$$

Any Questions?



Applying Toom-4 to $f(x)^2$, $f(x) = -1 - 2x - 3x^2 - 4x^3 + 4x^4 + 3x^5 + 2x^6 + x^7$

$$[(-1 - 2x) + (-3 - 4x)y + (4 + 3x)y^2 + (2 + x)y^3] = F(x, y) = G(x, y)$$



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$$[(-1 - 2x) + (-3 - 4x)y + (4 + 3x)y^2 + (2 + x)y^3] = F(x, y) = G(x, y)$$

y_0		$H(y_0) := F(x, y_0) \cdot G(x, y_0)$
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$$[(-1 - 2x) + (-3 - 4x)y + (4 + 3x)y^2 + (2 + x)y^3] = F(x, y) = G(x, y)$$

y_0		$H(y_0) := F(x, y_0) \cdot G(x, y_0)$
0	$H(0)$	$(-1 - 2x)^2 = 1 + 4x + 4x^2$
1	$H(1)$	$(2 - 2x)^2 = 4 - 8x + 4x^2$
2	$H(2)$	$(25 + 10x)^2 = 625 + 500x + 100x^2$
∞	$H(\infty)$	$(2 + x)^2 = 4 + 4x + x^2$
-1	$H(-1)$	$(4 + 4x)^2 = 16 + 32x + 16x^2$
-2	$H(-2)$	$(5 + 10x)^2 = 25 + 100x + 100x^2$
-3	$H(-3)$	$(-10 + 10x)^2 = 100 - 200x + 100x^2$



Applying Toom-4 (cont'd)

Write $H = c_0(x) + c_1(x)y + \dots + c_6(x)y^6$

$$\begin{bmatrix} c_0(x) \\ c_1(x) \\ c_2(x) \\ c_3(x) \\ c_4(x) \\ c_5(x) \\ c_6(x) \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 4 & 8 & 16 & 32 & 64 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & -1 & 1 & -1 & 1 & -1 & 1 \\ 1 & -2 & 4 & -8 & 16 & -32 & 64 \\ 1 & -3 & 9 & -27 & 81 & -243 & 729 \end{bmatrix}^{-1} \begin{bmatrix} H(0) \\ H(1) \\ H(2) \\ H(\infty) \\ H(-1) \\ H(-2) \\ H(-3) \end{bmatrix} = \begin{bmatrix} 1 + 4x + 4x^2 \\ 6 + 20x + 16x^2 \\ 1 + 2x + 4x^2 \\ -28 - 60x - 28x^2 \\ 4 + 2x + x^2 \\ 16 + 20x + 6x^2 \\ 4 + 4x + x^2 \end{bmatrix}$$

Applying Toom-4 (cont'd)

Write $H = c_0(x) + c_1(x)y + \dots + c_6(x)y^6$

$$\begin{bmatrix} c_0(x) \\ c_1(x) \\ c_2(x) \\ c_3(x) \\ c_4(x) \\ c_5(x) \\ c_6(x) \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 4 & 8 & 16 & 32 & 64 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & -1 & 1 & -1 & 1 & -1 & 1 \\ 1 & -2 & 4 & -8 & 16 & -32 & 64 \\ 1 & -3 & 9 & -27 & 81 & -243 & 729 \end{bmatrix}^{-1} \begin{bmatrix} H(0) \\ H(1) \\ H(2) \\ H(\infty) \\ H(-1) \\ H(-2) \\ H(-3) \end{bmatrix} = \begin{bmatrix} 1 + 4x + 4x^2 \\ 6 + 20x + 16x^2 \\ 1 + 2x + 4x^2 \\ -28 - 60x - 28x^2 \\ 4 + 2x + x^2 \\ 16 + 20x + 6x^2 \\ 4 + 4x + x^2 \end{bmatrix}$$

$$f(x)g(x) = 1 + 4x + (4 + 6)x^2 + 20x^3 + (16 + 1)x^4 + 2x^5 + (4 - 28)x^6 - 60x^7 + (-28 + 4)x^8 + 2x^9 + (1 + 16)x^{10} + 20x^{11} + (6 + 4)x^{12} + 4x^{13} + x^{14}$$