Advanced security notions for the SSH secure channel: theory and practice

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Based on joint work with Martin Albrecht, Jean Paul Degabriele and Torben Hansen

Information Security Group
Overview

1. Introducing SSH
2. SSH measurement study
3. An unfortunate sequence of attacks on CBC-mode in OpenSSH
4. Security analysis of other SSH and OpenSSH modes – CTR, ChaChaPoly, gEtM, AES-GCM.
5. Better security for SSH: InterMAC
6. Concluding remarks
Introducing SSH and related work
SecureShell or SSH is a network protocol that allows data to be exchanged using a secure channel between two networked devices. Used primarily on Linux and Unix based systems to access shell accounts, SSH was designed as a replacement for TELNET and other insecure remote shells, which send information, notably passwords, in plaintext, leaving them open for interception. The encryption used by SSH provides confidentiality and integrity of data over an insecure network, such as the Internet.

– Wikipedia
SSH Binary Packet Protocol

- Stateful Encode-then-E&M construction
- Packet length field measures the size of the packet: |PadLen| + |Payload| + |Padding|.
- RFC 4253 (2006): various block ciphers in CBC mode (with chained IV) and RC4.
Timeline of related work on SSH BPP

2002.

• Formal security analysis of SSH BPP by Bellare, Kohno and Namprempre [BKN02]: introduce stateful security notions for symmetric encryption and proved SSH-CTR and SSH-CBC variants (w/o IV chaining) secure.

2009.

• Albrecht, Paterson and Watson [APW09] discover a plaintext-recovery attack against SSH in CBC mode.

• The attack exploits fragmented delivery in TCP/IP, and works on all CBC variants considered in [BKN02].

• The then leading implementation was OpenSSH (reported 80% of servers); OpenSSH team release a patch in version 5.2 to stop the specific attack.
Timeline of related work on SSH BPP

2010.

• The [APW09] attack highlights deficiencies in the [BKN02] security model.
• Paterson and Watson [PW10] prove SSH-CTR secure in an extended security model that allows adversary to deliver fragmented ciphertexts.

2012.

• Boldyreva, Degabriele, Paterson and Stam [BDPS12] study ciphertext fragmentation more generally, addressing limitations in the [PW10] model, introducing IND-CFA security.
• [BDPS12] also considers boundary hiding and resistance to a special type of denial of service attack as additional security requirements.
SSH measurement study
In [ADHP16], we performed a measurement study of SSH deployment.

We conducted two complete IPv4 address space scans in Nov/Dec 2015 and Jan 2016 using ZGrab/Zmap.

- Grabbing banners and SSH servers’ preferred algorithms.
- Actual cipher used in a given SSH connection depends on client and server preferences.

Roughly $2^{24}$ servers found in each scan.

Nmap fingerprinting suggests mostly embedded routers and firewall devices.

Data available at:

https://bitbucket.org/malb/a-surfeit-of-ssh-cipher-suites/overview
## SSH versions

<table>
<thead>
<tr>
<th>software</th>
<th>scan 2015–12</th>
<th>scan 2016–01</th>
</tr>
</thead>
<tbody>
<tr>
<td>dropbear_2014.66</td>
<td>7,229,491</td>
<td>8,334,758</td>
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<tr>
<td>OpenSSH_5.3</td>
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</tr>
<tr>
<td>dropbear_2011.54</td>
<td>383,575</td>
<td>64,666</td>
</tr>
<tr>
<td>ROSSSH</td>
<td>345,916</td>
<td>333,992</td>
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<td>OpenSSH_5.1</td>
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<tr>
<td>OpenSSH_5.3p1</td>
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<td>0</td>
</tr>
<tr>
<td>OpenSSH_7.1</td>
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</table>

Mostly OpenSSH and dropbear; others less than 5%.
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<td>dropbear_2014.66</td>
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<td>(47.0%)</td>
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<td>56,082</td>
<td></td>
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<tr>
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<td>0</td>
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<tr>
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<td>83,793</td>
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Dropbear at 56-58%. 886k older than version 0.52, so vulnerable to variant of 2009 CBC-mode attack.
The state of SSH today: SSH versions

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OpenSSH at 37-39%. 166k older than version 5.2 and prefer CBC mode, so vulnerable to 2009 attack.
SSH versions

• Dropbear dominates over OpenSSH.

• Long tail of old software versions.
  • Most popular version of OpenSSH was version 5.3, released Oct 2009 (current version is 7.5).
  • Determined by major Linux distros?

• Non-negligible percentage of Dropbear and OpenSSH servers were potentially still vulnerable to the 2009 attack.
  • 8.4% for Dropbear.
### OpenSSH preferred algorithms

<table>
<thead>
<tr>
<th>Encryption and MAC Algorithm</th>
<th>Count</th>
<th>Percentage</th>
</tr>
</thead>
<tbody>
<tr>
<td>aes128-ctr + hmac-md5</td>
<td>3,877,790</td>
<td>(57.65%)</td>
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<tr>
<td>aes128-ctr + hmac-md5-ettm@</td>
<td>2,010,936</td>
<td>(29.90%)</td>
</tr>
<tr>
<td>aes128-ctr + umac-64-ettm@</td>
<td>331,014</td>
<td>(4.92%)</td>
</tr>
<tr>
<td>aes128-cbc + hmac-md5</td>
<td>161,624</td>
<td>(2.40%)</td>
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<td>chacha20-poly1305@</td>
<td>115,526</td>
<td>(1.72%)</td>
</tr>
<tr>
<td>aes128-ctr + hmac-shal</td>
<td>68,027</td>
<td>(1.01%)</td>
</tr>
<tr>
<td>des + hmac-md5</td>
<td>40,418</td>
<td>(0.60%)</td>
</tr>
<tr>
<td>aes256-gcm@</td>
<td>28,019</td>
<td>(0.42%)</td>
</tr>
<tr>
<td>aes256-ctr + hmac-sha2-512</td>
<td>17,897</td>
<td>(0.27%)</td>
</tr>
<tr>
<td>aes128-cbc + hmac-shal</td>
<td>11,082</td>
<td>(0.16%)</td>
</tr>
<tr>
<td>aes128-ctr + hmac-ripemd160</td>
<td>10,621</td>
<td>(0.16%)</td>
</tr>
</tbody>
</table>

OpenSSH preferred algorithms ("@"="@openssh.com")

- Lots of diversity (155 different combinations).
- CTR dominates, followed by CBC, surprising amount of EtM.
- ChaCha20-Poly1305 on the rise? (became default in OpenSSH 6.9).
- Small amount of GCM.
Dropbear preferred algorithms

- Less diversity than OpenSSH.
- CTR also dominates, followed by CBC.
- No “exotic” options.
- All CBC modes unpatched against variant of 2009 attack (8.4%).

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<thead>
<tr>
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<th>count</th>
</tr>
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<tbody>
<tr>
<td>aes128-ctr + hmac-shal-96</td>
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<tr>
<td>aes128-cbc + hmac-shal-96</td>
<td>478,181</td>
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<td>3des-cbc + hmac-shal</td>
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<tr>
<td>aes128-ctr + hmac-shal</td>
<td>62,465</td>
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<tr>
<td>aes128-ctr + hmac-sha2-256</td>
<td>36,150</td>
</tr>
<tr>
<td>aes128-cbc + hmac-shal</td>
<td>14,477</td>
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</tbody>
</table>
An unfortunate sequence of attacks on CBC mode in OpenSSH
SSH Binary Packet Protocol

How would you perform decryption for an incoming sequence of ciphertext fragments?
The [APW09] attack (simplified)

- Decryption in OpenSSH CBC mode (prior to 5.2):
  - Use a buffer to hold the incoming sequence of ciphertext fragments.
  - Decrypt the fragments block-by-block as they arrive.
  - 4-byte packet length field LF is obtained from the first block of the first fragment to be received.
  - Continue to buffer+decrypt until a total of LF+|MAC| bytes have been received.
  - Verify the MAC on SQN || PTXT (with connection termination and error message if MAC verification fails).
Breaking CBC mode in SSH [APWo9]

\[ C_{i-1}^* \rightarrow C_i^* \rightarrow d_K \rightarrow P_i^* \]

Target ciphertext block from stream
Breaking CBC mode in SSH [APW09]

$C_i^*$

Inject target block as first block of new ciphertext!
Breaking CBC mode in SSH [APWo9]

IV \rightarrow C_i^* \rightarrow d_K \rightarrow P_0' \rightarrow \text{Treated as length field}
Breaking CBC mode in SSH [APW09]
Breaking CBC mode in SSH [APW09]

- Once **enough** data has arrived, the receiver will get what it thinks is the MAC tag
  - The MAC verification will fail with overwhelming probability
  - So the connection is terminated (with an error message)
- **Question**: How much data is “enough” so that the receiver decides to check the MAC?
- Answer: whatever is specified in the length field:
Breaking CBC mode in SSH [APW09]

- Knowing IV and 32 bits of $P_0'$, the attacker can now recover 32 bits of the target plaintext block $P_i^*$.

$$\text{LF} \oplus [IV]_{0..3} = \oplus [C_{i-1}]_{0..3}$$
The [APW09] attack (less simplified)

• OpenSSH5.1 actually performs two sanity checks on the length field when decrypting the first ciphertext block:
  • Check 1: $5 \leq LF \leq 2^{18}$.
  • Check 2: total length ($LF+4$) is a multiple of the block size:
    $LF + 4 \mod BL = 0$.

• Each check produces a different error message on the network, distinguishable by attacker.

• If both checks pass, then OpenSSH waits for more bytes, then performs MAC check, resulting in a third distinct error message.

• The different error messages allow up to 32 bits of plaintext to be recovered with probability $2^{-18}$. 
OpenSSH 5.2 patch against [APW09] attack

Sanity checks:
5 ≤ LF ≤ $2^{18}$
LF + 4 mod BL = 0

→ PASS

Wait for LF+\(\mid MAC\) bytes

→ VERIFY

→ FAIL

ssh2_msg_disconnect

→ "corrupted MAC on input"

→ Wait until $2^{18}$ bytes have arrived, then check a MAC on $2^{18}$ bytes.

No error message is sent until $2^{18}$ bytes of ciphertext have arrived.

Is this a good patch?
OpenSSH 5.2 patch against [APW09] attack

Sanity checks:
\[ 5 \leq LF \leq 2^{18} \]
\[ LF + 4 \mod BL = 0 \]

\[
\begin{align*}
\text{PASS} & \quad \rightarrow \quad \text{FAIL} \\
\text{Wait for } LF + |\text{MAC}| \text{ bytes} & \quad \rightarrow \quad \text{FAIL} \\
\text{MAC on } \sim LF \text{ bytes + MAC on } 2^{18} \text{ bytes} & \quad \rightarrow \quad \text{WAIT until } 2^{18} \text{ bytes have arrived, then check a MAC on } 2^{18} \text{ bytes.}
\end{align*}
\]

No error message is ever sent until \( 2^{18} \) bytes of ciphertext have arrived.
[ADHP16] attack against the OpenSSH 5.2 patch

- Attacker can distinguish PASS/FAIL conditions, leaking 18 bits of plaintext.
- With careful timing, attacker can recover ~30 bits of plaintext.
OpenSSH 7.3 patch against [ADHP16] attack

Sanity checks:
- $5 \leq LF \leq 2^{18}$
- $LF + 4 \mod BL = 0$

PASS

Wait for $LF + |MAC|$ bytes

FAIL

$ssh2_msg_disconnect$

Wait until $2^{18}$ bytes have arrived, then check a MAC on $2^{18}$ bytes.

MAC on $2^{18}$ bytes

FAIL

"corrupted MAC on input"

Wait until $2^{18}$ bytes have arrived, then check a MAC on $2^{18}$ bytes.

Wait until $2^{18}$ bytes have arrived, then check a MAC on $2^{18} - LF$ bytes.

So is this a good patch?
Attacking the OpenSSH 7.3 patched patch

Wait a few seconds

Performed during the wait

Time

"Slow"

MAC on ~LF bytes +
MAC on $2^{18}$ - LF bytes
Sanity check PASS

"Fast"

MAC on $2^{18}$ bytes
Sanity check FAIL

Timing difference

1 byte

MAC error

$2^{18} - \text{BL} - 1$ bytes (quickly)
Attacking the OpenSSH 7.3 patched patch

Our recommended patch actually made things significantly worse!
I wonder if anyone noticed?

I think we got away with it!

I’m not so sure!
Disclosure of the attacks

• We first notified the OpenSSH team of the attack on the patch for the [APW09] attack on 5/5/2016.
• They first set of countermeasures in OpenSSH 7.3 (released 1/8/2016).
• We then notified OpenSSH of the new attack on 15/12/2016, along with some other, more subtle byte counting issues.
• These were partly addressed in OpenSSH 7.5 (released 20/3/2017).
• But several residual issues remain unpatched, including the final attack.
• In defence of OpenSSH:
  • OpenSSH has steadily been deprecating old algorithms and modes.
  • For example, CBC mode was already disabled by default in OpenSSH 6.7.
Security analysis of other SSH and OpenSSH modes – CTR, gEtM, AES-GCM, ChaCha20Poly1305
A number of new schemes have been introduced in OpenSSH since [APW09]:

- **AES-GCM**: since v6.2; **length field not encrypted** but is instead treated as associated data.

- **generic Encrypt-then-MAC (gEtM)**: since v6.2; overrides native E&M processing; **length field not encrypted** but protected by MAC.

- **ChaCha20-Poly1305@openssh.com**: since v6.5 and promoted to default in v6.9; **reintroduces encryption of length field**.
Binary Packet Protocol native E&M construction

**Diagram:**
- **Sequence Number:** 4
- **Packet Length:** 4
- **Pad Len:** 1
- **Payload**
- **Padding:** $\geq 4$
- **Encrypt**
- **Ciphertext**
- **MAC tag**

**Steps:**
1. Sequence Number
2. Packet Length
3. Pad Len
4. Payload
5. Padding
6. Encrypt
7. Ciphertext
8. MAC tag
Binary Packet Protocol generic EtM construction

- Stateful Encode-then-EtM construction.
- AES-GCM works similarly.
- Note packet length field in the clear.
- Code = documentation.
• Sequence: compute MAC, then decrypt, then check MAC.
• Issue arises because of retrofitting gEtM in legacy E&M code.
• No concrete attack, but dangerous to decrypt unauthenticated ciphertext (cf. padding oracle attacks).
• Addressed in OpenSSH 7.3.
ChaCha20-Poly1305@openssh.com

- ChaCha20-Poly1305@openssh.com: since OpenSSH 6.5 and promoted to default in v6.9; **reintroduces encryption of length field.**
- OpenSSH developers seem to care a lot about **hiding packet lengths!**
Security analysis from [ADHP16]

• We used the framework of [BDPS12] for symmetric encryption schemes supporting ciphertext fragmentation to analyse the security of these schemes.

• We identified and fixed a technical issue in the IND-sfCFA confidentiality definition from [BDPS12].

• We introduced a matching notion of ciphertext integrity, INT-sfCTXT, which was not considered in [BDPS12].
Security analysis from [ADHP16]

<table>
<thead>
<tr>
<th></th>
<th>IND-sfCFA</th>
<th>INT-sfCTF</th>
<th>BH-CPA</th>
<th>BH-sfCFA</th>
<th>n-DOS-sfCFA</th>
</tr>
</thead>
<tbody>
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<td>x</td>
<td>x</td>
</tr>
<tr>
<td>fixed-CBC</td>
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<td></td>
<td>✓</td>
<td>x</td>
<td>x</td>
</tr>
<tr>
<td>CTR</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>x</td>
<td>x</td>
</tr>
<tr>
<td>fgEtM</td>
<td>✓</td>
<td>✓</td>
<td></td>
<td>x</td>
<td>x</td>
</tr>
<tr>
<td>AES-GCM</td>
<td>✓</td>
<td>✓</td>
<td></td>
<td>x</td>
<td>x</td>
</tr>
<tr>
<td>ChaCha20-Poly1305</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>x</td>
<td>x</td>
</tr>
</tbody>
</table>

Security comparison of SSH AE modes

Additional goals from [BDPS12]:

- BH-CPA (passive adversary) – boundary hiding for passive attackers.
- BH-sfCFA (active adversary) – boundary hiding for active attackers.
- $n$-DOS-sfCFA: decryption must produce some output (plaintext or error) after receiving at most an $n$-bit sequence of fragments chosen by adversary.
InterMAC
An encryption scheme proposed in [BDPS12].
Parameterised by a positive integer $N$ (the chunk length).
Satisfies all 5 security notions:
- IND-sfCFA, IND-sfCTF, BH-CPA, BH-sfCFA, $(N + |MAC|)$-DOS-sfCFA.
Applies a generic EtM construction to chunks of data, incorporating additional metadata in the MAC computation.
Simple, easy to analyse construction; advanced security properties are intuitively obvious.
Small $N$: good DoS protection, but larger bandwidth overhead.
Idea: refine and implement InterMAC in OpenSSH to obtain stronger security than is currently available.
InterMAC

Chunk CTR
Msg CTR

Payload

N-1  N-1  N-1

"EtM"

$c_1$  $\tau_1$

$c_2$  $\tau_2$

$c_3$  $\tau_3$

$c_1$  $\tau_1$  $c_2$  $\tau_2$  $c_3$  $\tau_3$
• Use byte-oriented rather than bit-oriented format.
• Abandon underlying SSH packet format (so no length field, no padding byte, no random padding).
• Need some kind of plaintext padding (length not usually a multiple of N-1!): variant of ABYTE padding.
• Replace EtM with nonce-based AEAD, e.g. AES-GCM or ChaCha20-Poly1305.
• Chunk and message counter then become Associated Data, or are used to construct the nonce.
  • We choose the latter.
InterMAClib and OpenSSH

• C-implementation of InterMAC.
• Aim is to make the library easy to use for a developer.
• API: \texttt{im\_initialise}, \texttt{im\_encrypt}, \texttt{im\_decrypt}.
• Message counter and nonce management done by the library.
• Currently supports ChaCha-Poly and AES-GCM.
• Easy to extend with other AEAD schemes.
• POC integration into OpenSSH (v7.4).
• SSH InterMAC cipher suites: \texttt{im-aes128-gcm-N}, \texttt{im-chacha-poly-N}.
InterMAClib Throughput – SCP on Loopback

MB/s

- im-chacha-poly-4096
- im-aes128-gcm-4096
- im-chacha-poly-2048
- im-aes128-gcm-2048
- im-chacha-poly-1024
- im-aes128-gcm-1024
- im-chacha-poly-512
- im-aes128-gcm-512
- im-chacha-poly-256
- im-aes128-gcm-256
- im-chacha-poly-128
- im-aes128-gcm-128
- aes128-ctr + hmac-ripemd160
- aes128-cbc + hmac-sha1
- aes256-ctr + hmac-sha2-512
- aes128-gcm@
- 3des-cbc + hmac-md5
- aes128-ctr + hmac-sha1
- chacha20-poly1305@
- aes128-cbc + hmac-md5
- aes128-ctr + umac-64-etm@
- aes128-ctr + hmac-d5-etm@
- aes128-ctr + hmac-md5
InterMAClib Throughput – AWS_London to AWS_Oregon.
InterMAClib Total Bandwidth – AWS_London to AWS_Oregon.
Concluding Remarks
Concluding Remarks

• We have developed a deeper understanding of the diverse set of encryption modes available in (Open)SSH.
  • Measurement study, new attacks on CBC mode, security analysis
• None of the schemes in use possesses all the security properties desirable for SSH.
  • Boundary-hiding and DoS-resistance not achieved.
• Yet such schemes do exist, e.g. InterMAC from [BDPS12].
• In our on-going work, we are developing and prototyping efficient, provably secure alternatives that have all the desired properties.
Selected Literature


